Appl. No.: 09/945,106
Amdt. Dated: 11/01/2005
Allowed: 09/15/2005

Amendments to the Claims:

This listing of claims will replace all prior versions, and listings, of claims in the application:

Listing of Claims:

Claims 1-4 (canceled)

5. (previously presented): A method for differentiating congestion-related packet loss versus random packet loss in a wireless data connection, comprising:

monitoring changes in the length of a transmission queue in a wireless data connection;

designating packet loss as being due to congestion if said packet loss is preceded by an increase in the queue length;

designating packet loss as random loss if said packet loss is not preceded by an increase in the queue length;

monitoring changes in the length of said queue over an interval substantially equal to the amount of time it takes to transmit a window of data packets and receive acknowledgements corresponding to all data packets transmitted in the window;

initializing a state count to zero;

transitioning from the state count zero to state count one if the length of said queue increases during the next interval;

transitioning from the state count one to state count zero if the length of said queue decreases or remains steady during the next subsequent interval;

transitioning from state count one to state count two if the length of said queue increases during the next subsequent interval; and

designating packet loss as due to congestion if state count two is reached.

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6. (original): A method as recited in claim 5, further comprising:

applying a collision avoidance algorithm if packet loss is designated as due to congestion.

7. (original): A method as recited in claim 6, wherein said collision avoidance algorithm comprises reducing the sender's transmission window by one-half.

Claims 8-9 (canceled)

10. (previously presented): A method for differentiating congestion-related packet loss versus random packet loss in a wireless data connection, comprising:

monitoring changes in the length of a transmission queue in a wireless data connection;

designating packet loss as being due to congestion if said packet loss is preceded by an increase in the queue length;

designating packet loss as random loss if said packet loss is not preceded by an increase in the queue length;

determining whether congestion is developing in the forward or reverse path of the connection; and

isolating forward throughput from congestion on the reverse path;

wherein congestion is determined by calculating the relative delay that one packet experiences with respect to another as it traverses the connection.

11. (original): A method as recited in claim 10, wherein said relative delay is used to estimate the number of packets residing the in the queue.

12. (original): A method as recited in claim 11, further comprising keeping the number of packets in the queue at a minimum level by adjusting a congestion window.

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13. (original): A method as recited in claim 12, further comprising:

reducing the congestion window if the queue length increases; and

increasing the congestion window if the queue length decreases.

14. (currently amended): A TCP-based congestion management protocol for a wireless data connection, comprising:

monitoring changes in the length of a transmission queue in a data connection;

designating packet loss as being due to congestion if said packet loss is preceded by at least two consecutive intervals of increasing queue length; and

designating packet loss as random loss if said packet loss is not preceded by at least two consecutive intervals of increasing queue length;

applying a collision avoidance algorithm if packet loss is designated as due to congestion;

wherein each said interval comprises the amount of time it takes to transmit a window of data packets and receive acknowledgements corresponding to all data packets transmitted in the window;

wherein collision avoidance algorithm comprises

reducing the sender's transmission window by one-half;

initializing a state count to zero;

transitioning from a state count zero to ~~sate~~ state count one if the length of said queue increases during the next interval;

transitioning from the ~~sate~~ state count one to state count zero if the length of said queue decreases or remains steady during next subsequent interval;

transitioning from a state count one to state count two if the length of said queue increases during the next subsequent interval; and

designating packet loss as due to congestion if ~~sate~~ state count two is reached.

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Claims 15-20 (canceled)

21. (original): A protocol as recited in claim 14, further comprising determining whether congestion is developing in the forward or reverse path of the connection.

22. (original): A protocol as recited in claim 21, further comprising isolating forward throughput from congestion on the reverse path.

23. (original): A protocol as recited in claim 22, wherein congestion is determined by calculating the relative delay that one packet experiences with respect to another as it traverses the connection.

24. (original): A protocol as recited in claim 23, wherein said relative delay is used to estimate the number of packets residing the in the queue.

25. (original): A protocol as recited in claim 24, further comprising keeping the number of packets in the queue at a minimum level by adjusting a congestion window.

26. (original): A protocol as recited in claim 25, further comprising:
reducing the congestion window if the queue length increases; and
increasing the congestion window if the queue length decreases.

Claims 27-29 (canceled)

30. (previously presented): A method for differentiating congestion-related packet loss versus random packet loss in a wireless data connection, comprising:

monitoring changes in the length of a transmission queue in a wireless data connection over an interval substantially equal to the amount of time it takes to transmit a window of data packets and receive acknowledgements corresponding to all data packets transmitted in the window;

designating packet loss as being due to congestion if said packet loss is preceded by an increase in the queue length;

designating packet loss as random loss if said packet loss is not preceded by an increase in the queue length;

initializing a state count to zero;

transitioning from the state count zero to state count one if the length of said queue increases during the next interval;

transitioning from the state count one to state count zero if the length of said queue decreases or remains steady during the next subsequent interval;

transitioning from state count one to state count two if the length of said queue increases during the next subsequent interval; and

designating packet loss as due to congestion if state count two is reached.

31. (original): A method as recited in claim 30, further comprising:

applying a collision avoidance algorithm if packet loss is designated as due to congestion.

32. (original): A method as recited in claim 31, wherein said collision avoidance algorithm comprises reducing the sender's transmission window by one-half.

Claims 33-34 (canceled)

35. (previously presented): A method for differentiating congestion-related packet loss versus random packet loss in a wireless data connection, comprising:

monitoring changes in the length of a transmission queue in a wireless data connection over an interval substantially equal to the amount of time it takes to transmit a window of data packets and receive acknowledgements corresponding to all data packets transmitted in the window;

designating packet loss as being due to congestion if said packet loss is preceded by an increase in the queue length;

designating packet loss as random loss if said packet loss is not preceded by an increase in the queue length;

determining whether congestion is developing in the forward or reverse path of the connection; and

isolating forward throughput from congestion on the reverse path;

wherein congestion is determined by calculating the relative delay that one packet experiences with respect to another as it traverses the connection.

36. (original): A method as recited in claim 35, wherein said relative delay is used to estimate the number of packets residing the in the queue.

37. (original): A method as recited in claim 36, further comprising keeping the number of packets in the queue at a minimum level by adjusting a congestion window.

38. (original): A method as recited in claim 37, further comprising:

reducing the congestion window if the queue length increases; and

increasing the congestion window if the queue length decreases.

39. (currently amended): A TCP-based congestion management protocol for a wireless data connection, comprising:

monitoring changes in the length of a transmission queue in a data connection over an interval substantially equal to the amount of time it takes to transmit a window

of data packets and receive acknowledgements corresponding to all data packets transmitted in the window;

designating packet loss as being due to congestion if said packet loss is preceded by at least two consecutive intervals of increasing queue length; and

designating packet loss as random loss if said packet loss is not preceded by at least two consecutive intervals of increasing queue length;

applying a collision avoidance algorithm if packet loss is designated as due to congestion;

wherein each said interval comprises the amount of time it takes to transmit a window of data packets and receive acknowledgements corresponding to all data packets transmitted in the window;

wherein collision avoidance algorithm comprises

reducing the sender's transmission window by one-half;

initializing a state count to zero;

transitioning from a state count zero to ~~sate~~ state count one if the length of said queue increases during the next interval;

transitioning from the ~~sate~~ state count one to state count zero if the length of said queue decreases or remains steady during next subsequent interval;

transitioning from a state count one to state count two if the length of said queue increases during the next subsequent interval; and

designating packet loss as due to congestion if ~~sate~~ state count two is reached.

Claims 40-44 (canceled)

45. (original): A protocol as recited in claim 39, further comprising determining whether congestion is developing in the forward or reverse path of the connection.

46. (original): A protocol as recited in claim 45, further comprising isolating forward throughput from congestion on the reverse path.

47. (original): A protocol as recited in claim 46, wherein congestion is determined by calculating the relative delay that one packet experiences with respect to another as it traverses the connection.

48. (original): A protocol as recited in claim 47, wherein said relative delay is used to estimate the number of packets residing the in the queue.

49. (original): A protocol as recited in claim 48, further comprising keeping the number of packets in the queue at a minimum level by adjusting a congestion window.

50. (original): A protocol as recited in claim 49, further comprising:

reducing the congestion window if the queue length increases; and

increasing the congestion window if the queue length decreases.

51. (original): A method for improving TCP performance over a wireless connection, comprising:

detecting the initial stages of congestion in the connection, and identifying the direction of the congestion;

determining whether congestion is developing in the forward or reverse path of the connection;

isolating the forward throughput from events such as congestion that may occur on the reverse path;

determining congestion by calculating the relative delay that one packet experiences with respect to another as it traverses the network;

using said relative delay to estimate the number of packets residing in a

bottleneck queue;

keeping the number of packets in the bottleneck queue at a minimum level by adjusting the TCP source's congestion window;

reducing the congestion window if the bottleneck queue length increases; and

increasing the congestion window when the source detects additional bandwidth availability in the connection.